

## Lecture 8

①

SWAP test on multiple qubits  
↓ reducing circuit depth in  
the algorithm for increased  
circuit width.

suppose we have qubits in  
systems  $A_1, \dots, A_n \equiv A$

↓ in systems  $B_1, \dots, B_n \equiv B$

If we would like to swap

$A \leftrightarrow B$ , we can swap individually

That is,

$$\text{SWAP}_{AB} = \text{SWAP}_{A_1 B_1} \otimes \text{SWAP}_{A_2 B_2} \\ \otimes \dots \otimes \text{SWAP}_{A_n B_n}$$

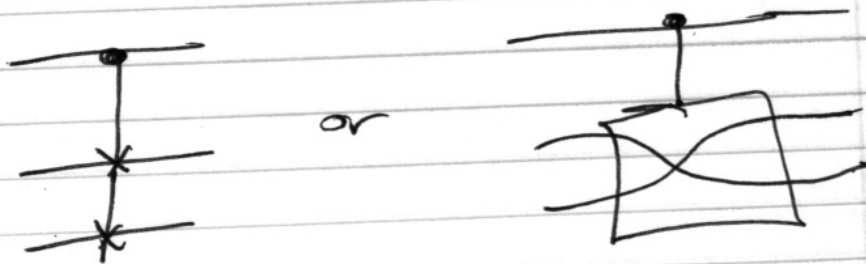
②

Now suppose that we have access to controlled-SWAP gate as a basic primitive (acting only on qubits)

This unitary is as follows:

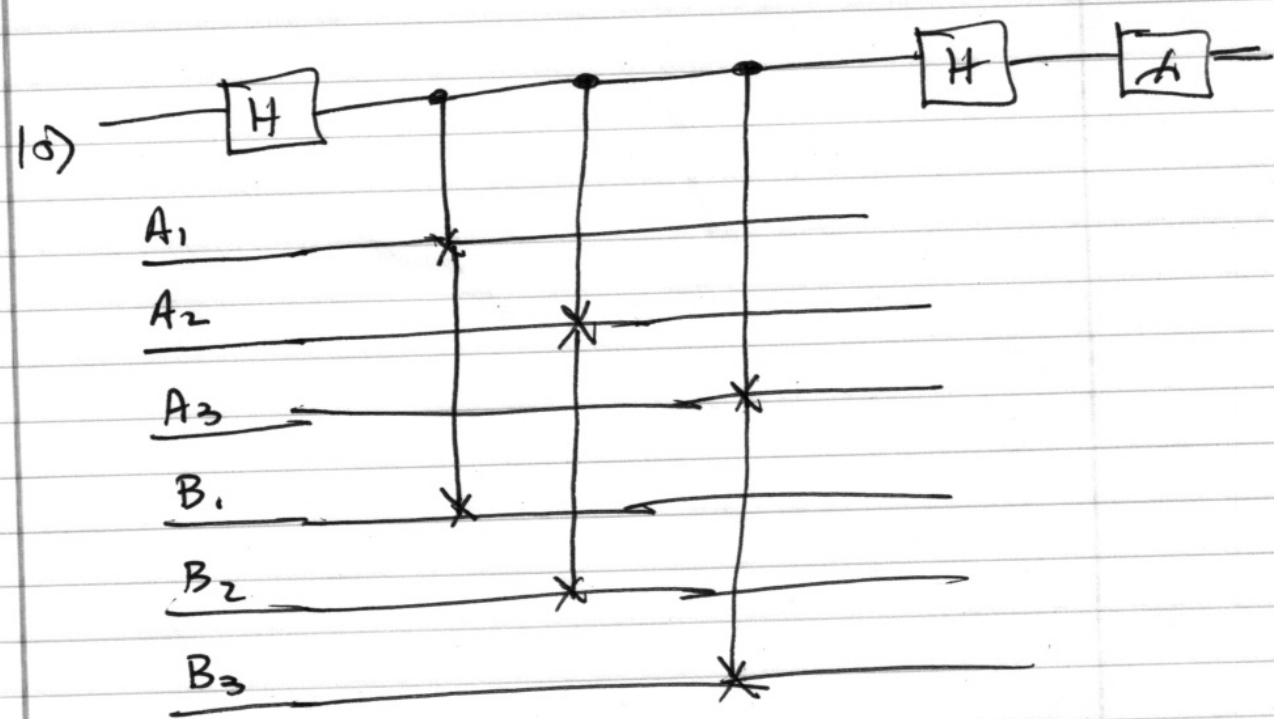
$$|s\rangle|t\rangle|c\rangle \otimes I_{A,B} \rightarrow |s\rangle|t\rangle|c\rangle \otimes \text{SWAP}_{A,B}$$

↓ written in circuit notation as



Then suppose we would like to implement the SWAP test for 3 qubits each. Then we can do it as follows:

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By similar reasoning as before,

$$\Pr[\sigma] = \frac{1 + |\langle \psi | \psi \rangle|^2}{2}$$

where  $|\psi\rangle_{A_1 A_2 A_3} \downarrow$

$|\psi\rangle_{B_1 B_2 B_3}$

$$\downarrow \Pr[\downarrow] = \frac{1 - |\langle \psi | \psi \rangle|^2}{2}$$

Here we see that the circuit depth is linear in the # of qubits. Is

there a way to reduce the circuit depth?

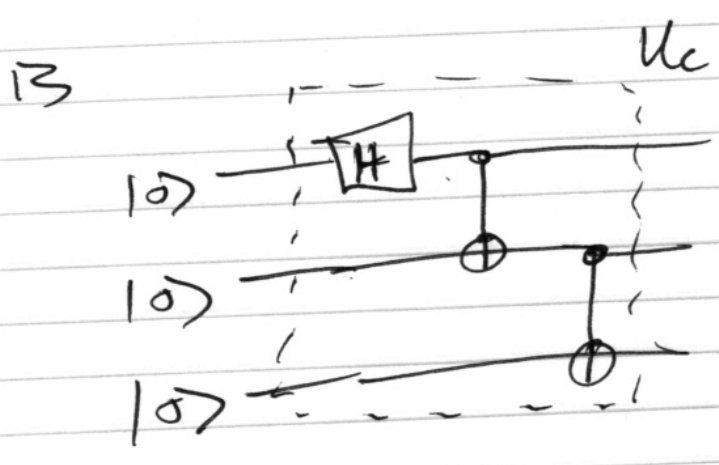
Consider that the bottleneck is that the controls of the SWAP gates all act on the same control qubit.

We can make this part of the circuit have constant depth by a technique we can call q. fan-out

Consider instead preparing  
 a GHZ state of three  
 qubits & having three  
 control qubits instead of  
 a single one.

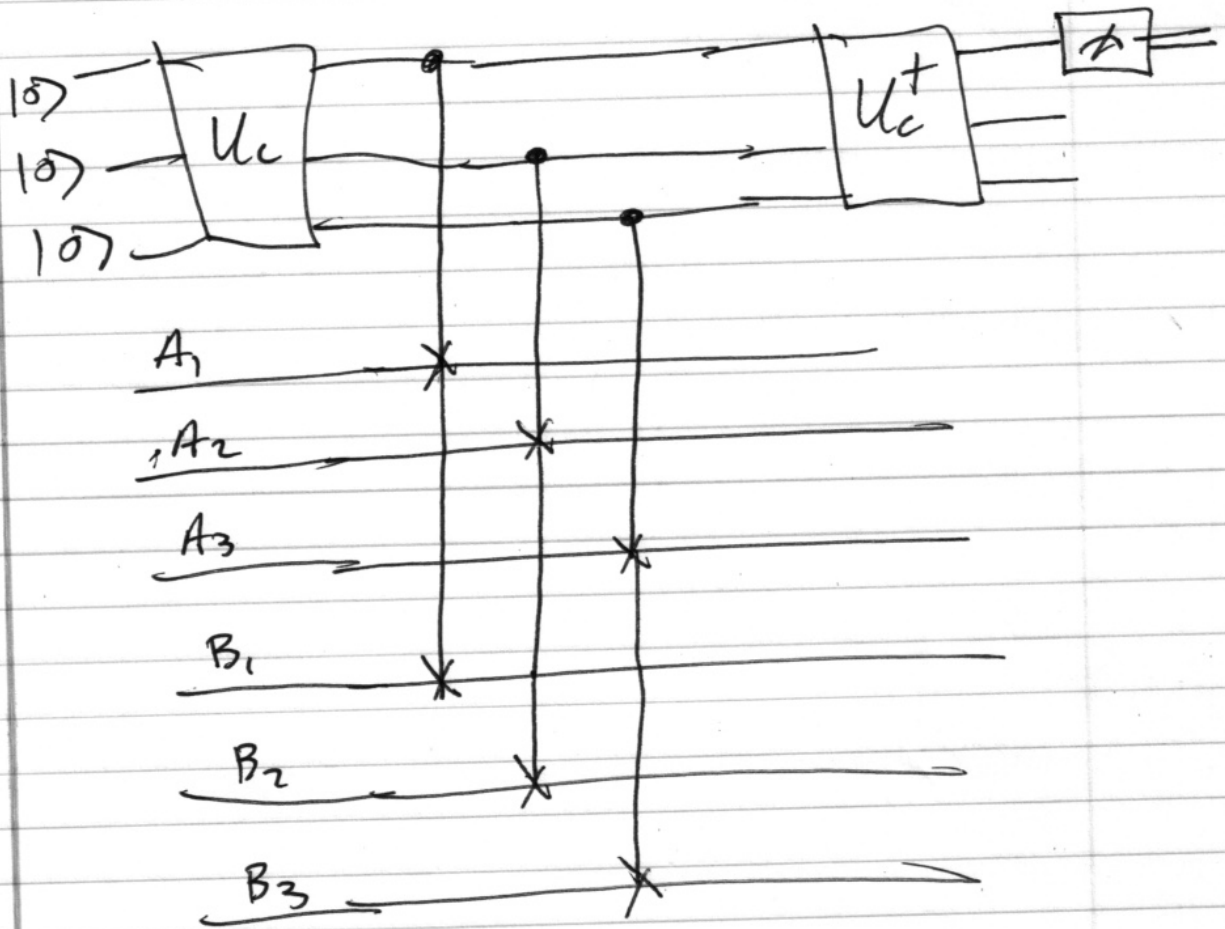
circuit to encode a GHZ state

$$\frac{|000\rangle + |111\rangle}{\sqrt{2}}$$



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then we can modify the circuit to be  
a  
SWAP test



these gates can be

run in parallel

∴ so this part now has

constant depth.

this implies that this part of the algorithm  
will be faster & more fewer errors

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let us now analyze the algorithm

Initial state is

$$|000\rangle |\psi\rangle_{A_1 A_2 A_3} |\varphi\rangle_{B_1 B_2 B_3}$$

After  $U_c$

$$\rightarrow \frac{|000\rangle + |111\rangle}{\sqrt{2}} |\psi\rangle_A |\varphi\rangle_B$$

$$= \frac{1}{\sqrt{2}} \cancel{|000\rangle + |111\rangle}$$

$$\frac{1}{\sqrt{2}} (|000\rangle |\psi\rangle_A |\varphi\rangle_B + |111\rangle |\varphi\rangle_A |\psi\rangle_B)$$

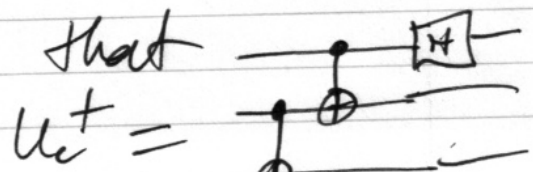
After controlled - SWAPs,

keeping in mind the local actions,

$$\rightarrow \frac{1}{\sqrt{2}} (|000\rangle |\psi\rangle_A |\varphi\rangle_B + |111\rangle |\varphi\rangle_A |\psi\rangle_B)$$

Now apply  $U_c^\dagger$  to control qubits,

keeping in mind that



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After the CNOTs, the state becomes

$$\begin{aligned} &\rightarrow \frac{1}{\sqrt{2}} \left( |000\rangle |\psi\rangle_A |\psi\rangle_B + |100\rangle |\psi\rangle_A |\psi\rangle_B \right) \\ &= \frac{1}{\sqrt{2}} \left( |0\rangle |\psi\rangle_A |\psi\rangle_B + |1\rangle |\psi\rangle_A |\psi\rangle_B \right) |00\rangle \end{aligned}$$

↑  
These qubits  
get  
factored out

then rest of analysis proceeds as before & we conclude

$$\text{that } \Pr[\sigma] = \frac{1 + |\langle \psi | \psi \rangle|^2}{2}$$

$$\Pr[\downarrow] = \frac{1 - |\langle \psi | \psi \rangle|^2}{2}$$

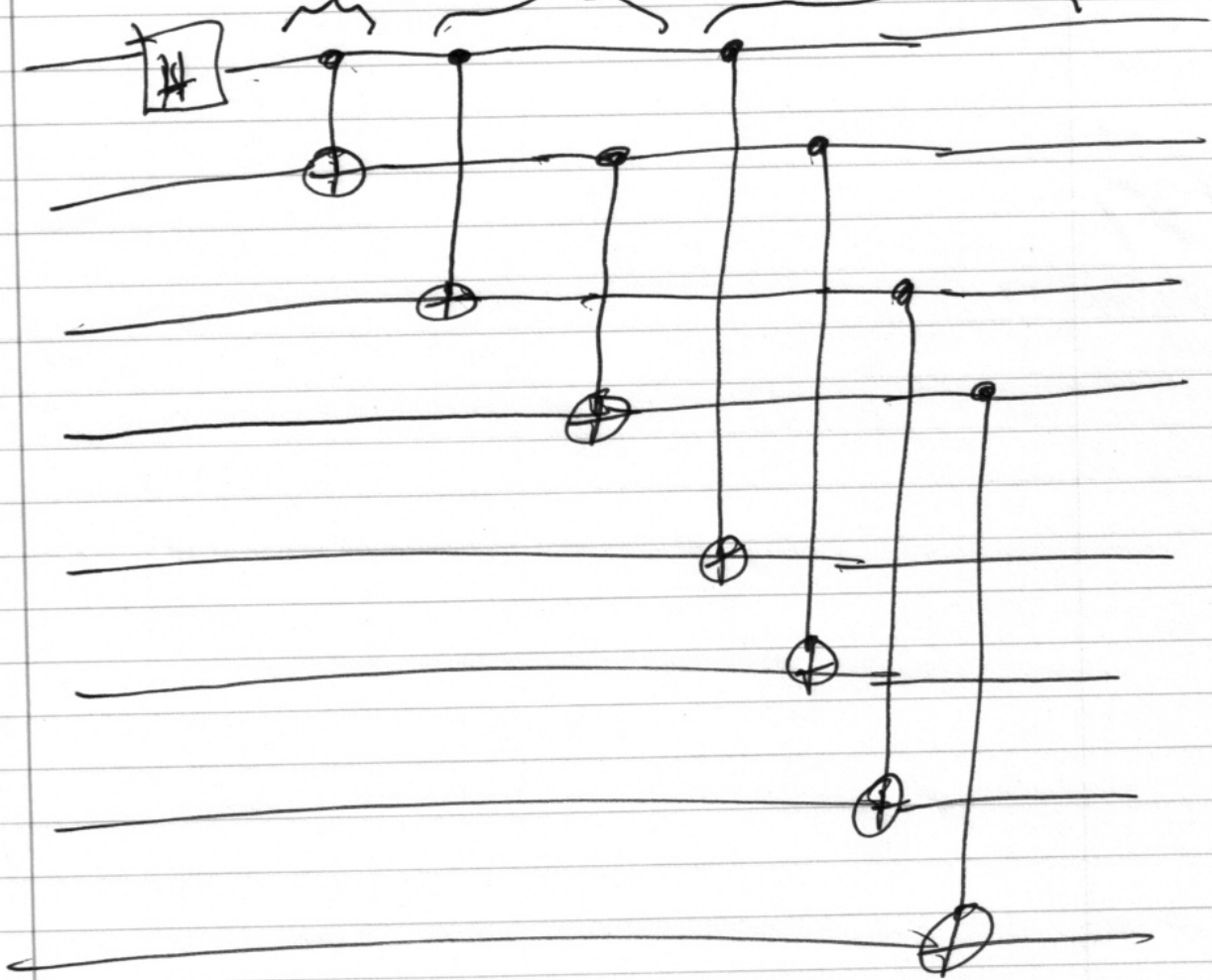


The circuit depth, as of now,

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~~there is a~~ for preparing  
the GHZ state, is  
logarithmic depth.

General construction is  
1 layer 1 layer 1 layer

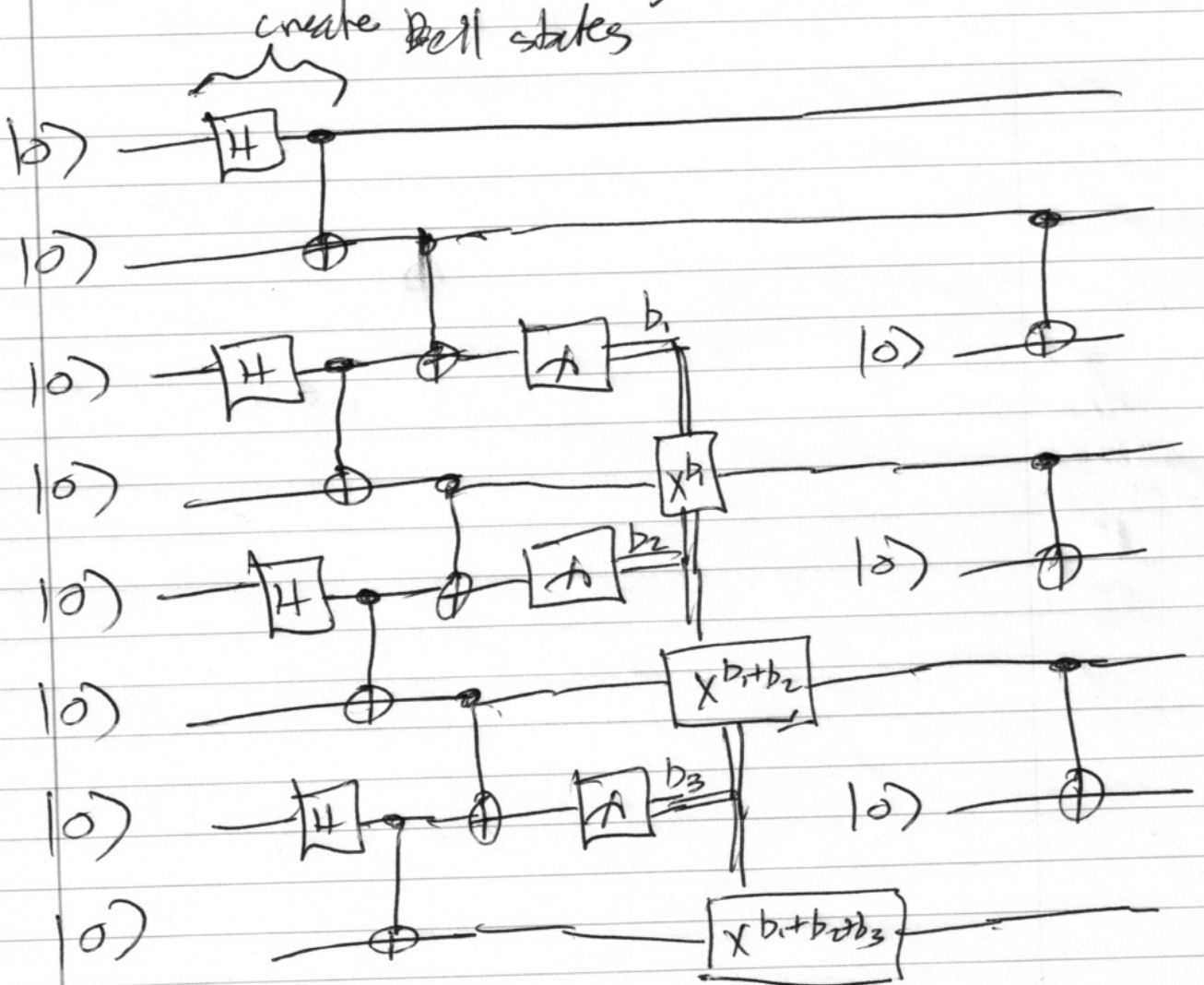


Can we do better?

Yes!

We can prepare GHZ states in constant quantum depth w/ ~~at~~ mid-circuit measurements

Consider the following circuit



This is a constant quantum depth circuit

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Why does this work?

Consider a simple example

$$|0\rangle|0\rangle|0\rangle|0\rangle$$

$$\rightarrow |\Phi^+\rangle|\Phi^+\rangle$$

$$= \frac{1}{\sqrt{2}}(|00\rangle+|11\rangle) \otimes \frac{1}{\sqrt{2}}(|00\rangle+|11\rangle)$$

$$= \frac{1}{2}(|0000\rangle+|1100\rangle+|0011\rangle+|1111\rangle)$$

after CNOT

$$\rightarrow \frac{1}{2}(|0000\rangle+|1110\rangle+|0011\rangle+|1101\rangle)$$

measuring the third qubit

gives

$$\frac{1}{\sqrt{2}}(|000\rangle+|111\rangle) \text{ w/ prob. } 1/2$$

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$$\frac{1}{\sqrt{2}} \left( |1110\rangle + |0011\rangle \right) \text{ w/ prob. } \frac{1}{2}$$

if second happens, then apply

X to third qubit

to get

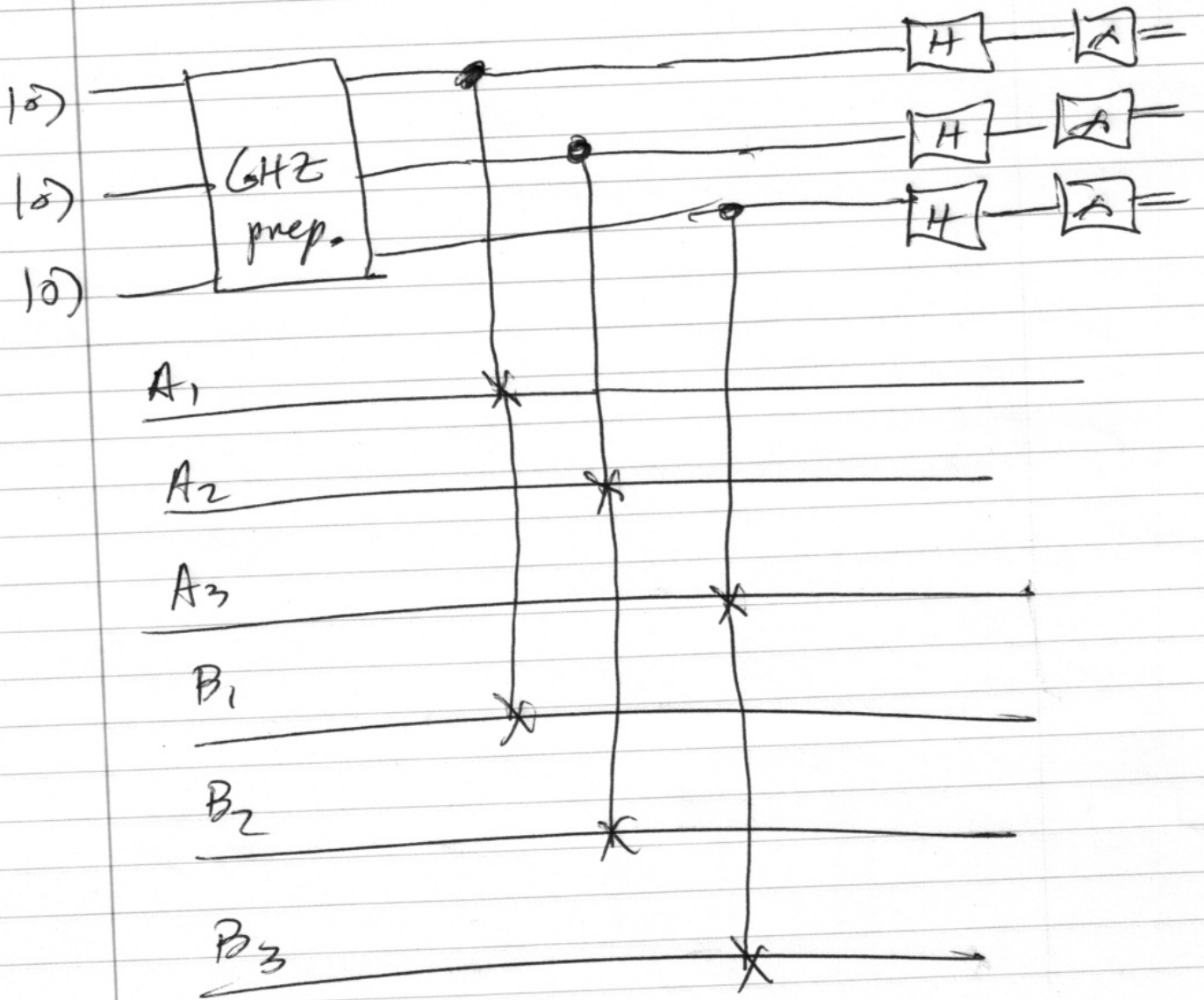
$$\frac{1}{\sqrt{2}} \left( |1000\rangle + |1111\rangle \right)$$

circuit provided generalizes  
this construction

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so now the preparation is  
constant depth and so  
are controlled SWAPs.  
What about final part?

Modifies to be as follows:



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As before, after controlled SWAPs,  
state is

$$\frac{1}{\sqrt{2}} \left( |000\rangle |\psi\rangle |\psi\rangle + |111\rangle |\psi\rangle |\psi\rangle \right)$$

what happens after measuring  
in Hadamard basis?

$$\text{Define } |\tilde{x}\rangle = \frac{1}{\sqrt{2}} \left( |0\rangle + (-1)^x |1\rangle \right)$$

for  $x \in \{0, 1\}$

probability of getting outcome

$x_1 x_2 x_3$  is

$$\left\| \langle \tilde{x}_1 \tilde{x}_2 \tilde{x}_3 | \frac{1}{\sqrt{2}} \left( |000\rangle |\psi\rangle |\psi\rangle + |111\rangle |\psi\rangle |\psi\rangle \right) \right\|^2$$

$$= \left\| \frac{1}{\sqrt{2^3}} \left( |\psi\rangle|\varphi\rangle + (-1)^{x_1+x_2+x_3} |\varphi\rangle|\psi\rangle \right) \right\|_2^2$$

$$= \frac{1}{2^3} \left( 2 + (-1)^{x_1+x_2+x_3} 2 \cancel{2} |\langle\psi|\varphi\rangle|^2 \right)$$

$$= \frac{1}{2^3} \left( 1 + (-1)^{x_1+x_2+x_3} |\langle\psi|\varphi\rangle|^2 \right)$$

Now set a R.V.

$$Z = (-1)^{x_1+x_2+x_3}$$

after observing measurement outcomes  $x_1, x_2, x_3$ .

Then

$$\mathbb{E}[Z] = \sum_{x_1, x_2, x_3} \text{Pr}[x_1, x_2, x_3] (-1)^{x_1+x_2+x_3}$$

$$= \sum_{x_1, x_2, x_3} \frac{1}{2^3} \left( 1 + (-1)^{x_1+x_2+x_3} |\langle\psi|\varphi\rangle|^2 \right) \cdot (-1)^{x_1+x_2+x_3}$$

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$$= \frac{1}{2^3} \sum_{x_1, x_2, x_3} \left[ (-1)^{x_1 + x_2 + x_3} + |\langle \psi | \psi \rangle|^2 \right]$$

$$= |\langle \psi | \psi \rangle|^2 \Rightarrow \mathbb{E}[z] = |\langle \psi | \psi \rangle|^2$$

So this gives a constant  
depth circuit to  
evaluate  $|\langle \psi | \psi \rangle|^2$