

Lecture 17.2

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Q. Eigenvalue transforms:

- previously, we discussed transformation of a single matrix element.
- Now, extend to transform all eigenvalues of a Hermitian matrix H (Hamiltonian) embedded in a unitary.
(Sps. $H \in \mathbb{R}^{N \times N}$)

Suppose now that we have this unitary U that block encodes H as

$$U = \begin{bmatrix} H & \sqrt{I - H^2} \\ \sqrt{I - H^2} & -H \end{bmatrix}$$

$$U = \sigma_z \otimes H + \sigma_x \otimes \sqrt{I - H^2} \quad (2)$$

Suppose that H has an eigendecomposition

$$\text{as } H = \sum_{\lambda} \lambda |\lambda\rangle\langle\lambda|$$

$$\text{Then } \sqrt{I - H^2} = \sum_{\lambda} \sqrt{1 - \lambda^2} |\lambda\rangle\langle\lambda|$$

\Rightarrow

$$U = \sigma_z \otimes H + \sigma_x \otimes \sqrt{I - H^2}$$

$$= \sigma_z \otimes \sum_{\lambda} \lambda |\lambda\rangle\langle\lambda| + \sigma_x \otimes \sum_{\lambda} \sqrt{1 - \lambda^2} |\lambda\rangle\langle\lambda|$$

$$= \sum_{\lambda} \lambda \sigma_z \otimes |\lambda\rangle\langle\lambda| + \sqrt{1 - \lambda^2} \sigma_x \otimes |\lambda\rangle\langle\lambda|$$

$$= \sum_{\lambda} (\lambda \sigma_z + \sqrt{1 - \lambda^2} \sigma_x) \otimes |\lambda\rangle\langle\lambda|$$

$$= \sum_{\lambda} R(\lambda) \otimes |\lambda\rangle\langle\lambda|$$

$R(\lambda)$ is ~~matrix~~ matrix from before.

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If we switch ordering of tensor product to be as follows, then we get direct sum

$$\sum_k |k\rangle\langle k| \otimes R(k)$$

$$= \bigoplus_k R(k)$$

We thus have N Bloch spheres & we can rotate them all simultaneously by using a common rotation operator.

Alternatively, we have uncovered N signal unitaries ($\{R(k)\}_{k=1}^N$) that can be processed simultaneously by some signal processing unitaries

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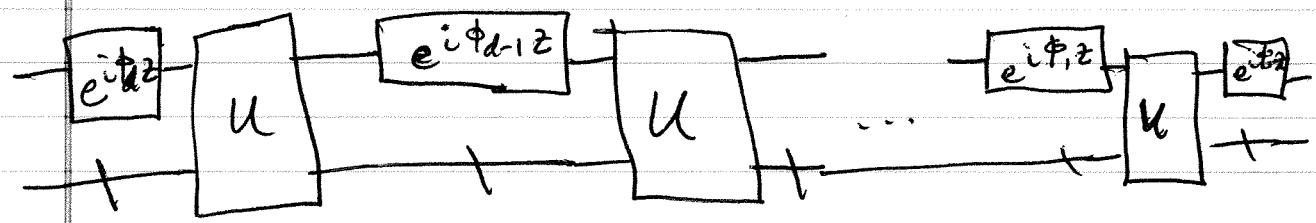
- For the special form of U discussed, we can process all N signals $(\sum \xi_r)$ simultaneously by using signal processing unitaries of the form $e^{i\phi z} = S(\phi)$.

- That is, we would do

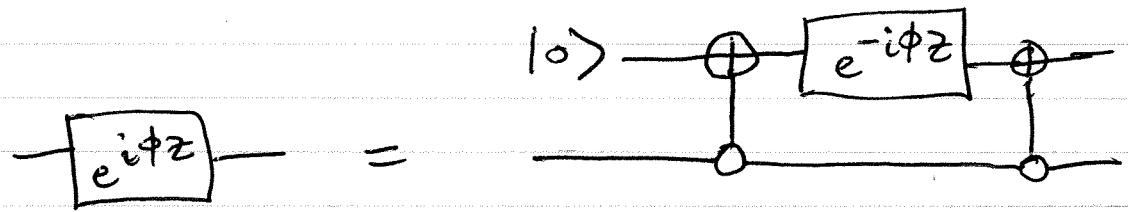
$$\begin{aligned} & (e^{i\phi_0 z} \otimes I) \prod_{k=1}^d U(e^{i\phi_k z} \otimes I) \\ &= (e^{i\phi_0 z} \otimes I) \prod_{k=1}^d \left(\sum_{\lambda} R(\lambda) \otimes |\lambda\rangle\langle\lambda| \right) e^{i\phi_k z} \otimes I \\ &= \sum_{\lambda} e^{i\phi_0 z} \underbrace{\left(\prod_{k=1}^d R(\lambda) e^{i\phi_k z} \right)}_{\text{signal processing of all } N \text{ signals simultaneously}} \otimes |\lambda\rangle\langle\lambda| \end{aligned}$$

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Thus, the circuit looks like



An alternative way for realizing rotations is by making use of observation that



Indeed,
$$e^{i\phi z} |\psi\rangle = e^{i\phi z} (\alpha |0\rangle + \beta |1\rangle)$$

$$= \alpha e^{i\phi} |0\rangle + \beta e^{-i\phi} |1\rangle$$

second circuit gives

$$|\psi\rangle |0\rangle = \alpha |100\rangle + \beta |110\rangle$$

CNOT

$$\rightarrow \alpha |101\rangle + \beta |110\rangle$$

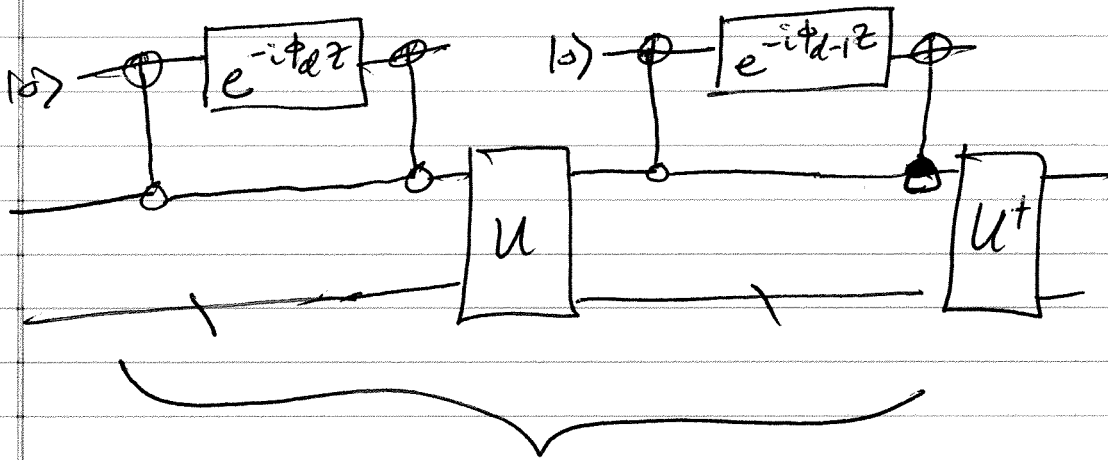
z-rot

$$\rightarrow \alpha e^{i\phi} |101\rangle + \beta e^{-i\phi} |110\rangle$$

$$\rightarrow (\alpha e^{i\phi} |10\rangle + \beta e^{-i\phi} |11\rangle) |0\rangle$$

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In this case, observe that $U = U^\dagger$, so that circuit can be written as



repeat $d/2$ more times
(assuming d)

Now identifying $\Pi = |0\rangle\langle 0| \otimes I$ (d even)

we can write

$$e^{i\phi\sigma_z} \otimes I = e^{i\phi(|0\rangle\langle 0| - |1\rangle\langle 1|)} \otimes I$$

$$= e^{i\phi(|0\rangle\langle 0| \otimes I - |1\rangle\langle 1| \otimes I)}$$

$$= e^{i\phi(\Pi - (I - \Pi))}$$

$$= e^{i\phi(2\Pi - I)}$$

this form is useful for generalizations

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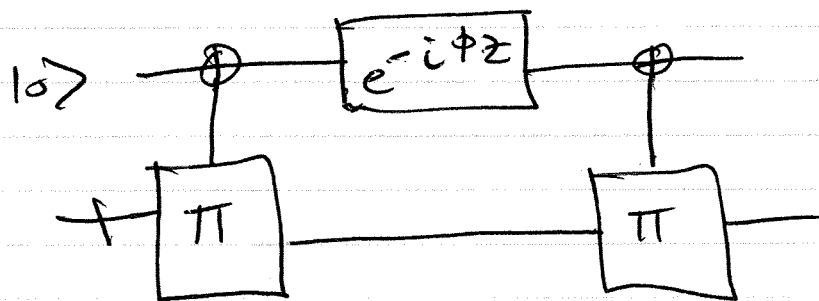
in which ~~U~~^U does not have
a special form. Suppose
instead that

$$U = \begin{bmatrix} H & \cdot \\ \cdot & \cdot \end{bmatrix}$$

where other entries are arbitrary.

Now let Π denote the projection
onto the top left block.

Then $e^{i\phi(2\Pi - I)} = \Pi e^{i\phi} \Pi$
is the
appropriate generalization &
can be realized by



where Π -controlled-Not is

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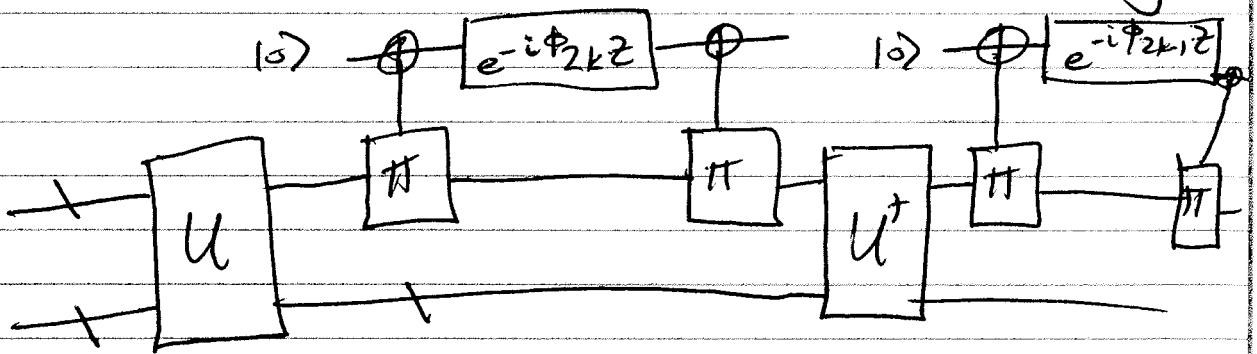
$$\pi \text{CNOT} = \pi \otimes X + (I - \pi) \otimes I$$

Flip the ~~second~~ ^{last} qubit if

1st register is in subspace onto which π projects

Most general ^{unitary} circuit for even d looks like

$$\prod_{k=1}^{d/2} \pi_{\phi_{2k-1}} U^\dagger \pi_{\phi_{2k}} U = \begin{bmatrix} \text{Poly}(H) & \cdot \\ \cdot & \cdot \\ \cdot & \cdot \end{bmatrix}$$



repeat $d/2$ times.

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Can then use polynomial transformations
for a variety of purposes

One example: eigenvalue filtering

- take eigenvalues above/below
a threshold to one &
others to zero,

q. singular value transformation

q. signal processing sequences
can be used to transform
all the singular values of
a rectangular matrix A .

Recall SVD of a matrix A :

$$A = W \Sigma V^+$$

where W & V are unitaries

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σ Σ is a diagonal, rectangular matrix of singular values.

denote columns of U by $\{|u_k\rangle\}_k$
& " " V by $\{|v_k\rangle\}_k$
each orthonormal bases.

can rewrite SVD as

$$A = \sum_{k=1}^n \sigma_k |u_k\rangle \langle v_k|$$

Given Unitary U such that
 A is block encoded as

$$U = \tilde{\Pi} \begin{bmatrix} A & \cdot \\ \cdot & \cdot \end{bmatrix} \Pi$$

where $\tilde{\Pi} = \sum_k |u_k\rangle \langle u_k|$ &
 $\Pi = \sum_k |v_k\rangle \langle v_k|$

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these are projectors that
locate A w/in U as

$$A = \tilde{\Pi} U \Pi$$

Going forward, for simplicity,
let us suppose that A is
square & U has form

$$U = \begin{bmatrix} A & \sqrt{I - A^2} \\ \sqrt{I - A^2} & -A \end{bmatrix}$$

$$\sqrt{I - A^2} = \sum_k \sqrt{1 - \sigma_k^2} |w_k\rangle \langle v_k|$$

can verify that $U^\dagger U = I$

can again write U as

$$U = \sigma_z \otimes A + \sigma_x \otimes \sqrt{I - A^2}$$

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$$= \sigma_z \otimes \sum_k \sigma_k \otimes |w_k\rangle\langle v_k| +$$

$$\sigma_x \otimes \sum_k \sqrt{1-\sigma_k^2} \otimes |w_k\rangle\langle v_k|$$

$$= \sum_k \left(\sigma_k \sigma_z + \sqrt{1-\sigma_k^2} \sigma_x \right) \otimes |w_k\rangle\langle v_k|$$

$$= \sum_k R(\sigma_k) \otimes |w_k\rangle\langle v_k|$$

↓

$$U^\dagger = \sum_k R(\sigma_k) \otimes |v_k\rangle\langle w_k|$$

can then realize a q. signal processing of singular values as

$$(e^{i\phi_0 z} \otimes I) U \prod_{k=1}^{(l-1)/2} (e^{i\phi_{2k} z} \otimes I) U^\dagger (e^{i\phi_{2k+1} z} \otimes I) U$$

$$= \begin{bmatrix} \text{Poly}^{(sv)}(A) & \cdot \\ \cdot & \cdot \end{bmatrix}$$

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where $Poly^{(SU)}(A)$

$$= \sum_k Poly(\sigma_k) |w_k\rangle\langle v_k|$$

More generally, transformation

for odd $d \geq 3$

$$\tilde{\Pi}_{\phi_1} U \left[\prod_{k=1}^{(d-1)/2} \Pi_{\phi_{2k}} U^\dagger \tilde{\Pi}_{\phi_{2k+1}} U \right]$$

Main difference between QSUT &

eigenvalue transformation is

that the ^{Bloch sphere} transformations of

U switch between bases $\{|v_k\rangle\}_k$

& $\{|w_k\rangle\}_k$

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How to block encode?

Many applications assume that

A is available ~~as~~ encoded as
a block in U .

To gain some intuition, suppose that

- Unitary A is available,

then controlled- A is a block
encoding of A .

- Specifically,

$$|0\rangle\langle 0| \otimes A + |1\rangle\langle 1| \otimes I$$

$$= \begin{bmatrix} A & 0 \\ 0 & I \end{bmatrix}, \quad \forall \text{ so } A \text{ is}$$

available as block encoded operators.

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If controlled A is not available,
but \exists eigenvector $|x\rangle$ of A ,
can then do

