

Lecture 10

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three noiseless protocols

- entanglement distribution
- super-dense coding
- teleportation

- resource counting
- proofs of optimality

Nonlocal unit resources

1) noiseless qubit channel

$$|i\rangle^A \rightarrow |i\rangle^B$$

preserves superpositions so that

$$\alpha|0\rangle^A + \beta|1\rangle^A \rightarrow \alpha|0\rangle^B + \beta|1\rangle^B$$

indicate this resource w/ notation:

$$[q \rightarrow q]$$

2) Noiseless classical bit channel

$$|i\rangle\langle i|^A \rightarrow |i\rangle\langle i|^B$$

$$\neq |i\rangle\langle j| \rightarrow 0 \quad \text{for } i \neq j$$

resource is

$$[c \rightarrow c]$$

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completely dephasing channel
is example

$$\begin{aligned} \Lambda(\rho) &\rightarrow \frac{1}{2} \rho + \frac{1}{2} Z \rho Z \\ &= |0\rangle\langle 0| \rho |0\rangle\langle 0| + |1\rangle\langle 1| \rho |1\rangle\langle 1| \\ &= \langle 0| \rho |0\rangle |0\rangle\langle 0| + \langle 1| \rho |1\rangle |1\rangle\langle 1| \end{aligned}$$

preserves a classical mixture

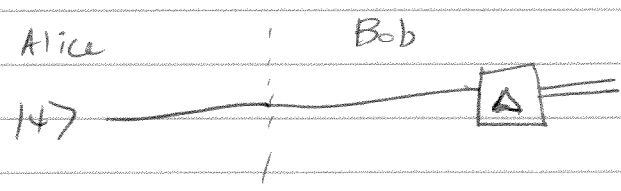
$$\frac{1}{2} (|0\rangle\langle 0| + |1\rangle\langle 1|)$$

but destroys a superposition state

$$\begin{aligned} |+\rangle\langle +| &= \left(\frac{|0\rangle + |1\rangle}{\sqrt{2}} \right) \left(\frac{\langle 0| + \langle 1|}{\sqrt{2}} \right) = \frac{1}{2} (|0\rangle\langle 0| + |0\rangle\langle 1| + |1\rangle\langle 0| + |1\rangle\langle 1|) \\ &\rightarrow \frac{1}{2} (|0\rangle\langle 0| + |1\rangle\langle 1|) \end{aligned}$$

similar to what measurement does

noiseless qubit channel can simulate
noiseless classical:



Thus, $\{q \rightarrow q\} \geq \{c \rightarrow c\}$

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Ebit: shared nonlocal entanglement

$$\frac{1}{\sqrt{2}} \left(|0\rangle^A |0\rangle^B + |1\rangle^A |1\rangle^B \right)$$

Classical comm. alone cannot simulate entanglement or quantum comm.

Entanglement alone cannot simulate classical or quantum comm.

1st Protocol: Entanglement Distribution

Two steps

- 1) create Bell state locally
- 2) send half of it over a noiseless qubit channel to create e-bit

Diagram



$$[a \rightarrow a] \rightarrow [eq]$$

- 1) becomes ebit
- 2) only count nonlocal resource
- 3) operations are perfect

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could we have

$$[qq] \geq [q \rightarrow q] \quad ?$$

no, violates no-signaling

\therefore quantum comm. is a stronger comm. resource than entanglement

2nd protocol: Super-dense coding

But first, elementary coding

- 1) Alice prepares $|0\rangle$ or $|1\rangle$
- 2) transmits over noiseless qubit channel
- 3) Bob measures Z

$$\therefore [q \rightarrow q] \geq [c \rightarrow c]$$

cannot do better than this

(homework exercise +
more generally
Holevo bound)

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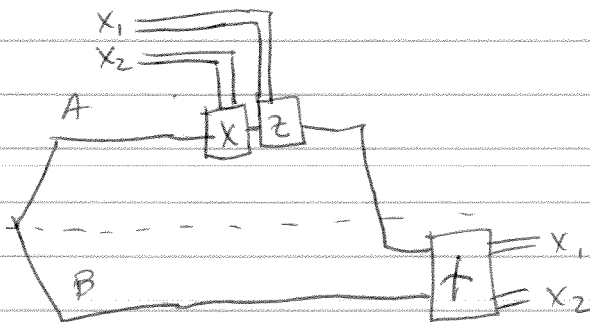
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dense coding

- 1) Alice + Bob share an ebit $|\Phi^+\rangle^{AB}$.
Alice applies a Pauli operator to her half of the ebit
one of $\{I, X, Z, XZ\}$. state becomes one of

$$|\Phi^+\rangle^{AB}, |\Phi^-\rangle^{AB}, |\Psi^+\rangle^{AB}, |\Psi^-\rangle^{AB}$$

- 2) transmits A over noiseless qubit channel
- 3) Bob performs Bell measurement to distinguish these states



resource inequality is

$$[qq] + [q \rightarrow q] \geq 2[c \rightarrow c]$$

Note: could have done

Also: protocol is private

$$2[q \rightarrow q] \geq 2[c \rightarrow c]$$

but dense coding only requires weaker resource of entanglement

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3rd protocol : teleportation

Alice & Bob "swap their equipment"

Alice & Bob share ebit $|\Phi^+\rangle_{AB}$

Alice wants to transmit qubit

$$|\psi\rangle^{A'} \equiv \alpha|0\rangle^{A'} + \beta|1\rangle^{A'}$$

overall state is

$$\begin{aligned} & |\psi\rangle^{A'} |\Phi^+\rangle_{AB} \\ &= (\alpha|0\rangle^{A'} + \beta|1\rangle^{A'}) \left(\frac{|100\rangle_{AB} + |111\rangle_{AB}}{\sqrt{2}} \right) \\ &= \frac{1}{\sqrt{2}} \left(\alpha|1000\rangle^{A'AB} + \alpha|1011\rangle^{A'AB} + \right. \\ & \quad \left. \beta|1100\rangle^{A'AB} + \beta|1111\rangle^{A'AB} \right) \end{aligned}$$

use results from homework to write as

$$\begin{aligned} &= \frac{1}{2} \left\{ \alpha \left(|\Phi^+\rangle^{A'A} + |\Phi^-\rangle^{A'A} \right) |0\rangle^B + \right. \\ & \quad \beta \left(|\Phi^+\rangle^{A'A} - |\Phi^-\rangle^{A'A} \right) |0\rangle^B + \\ & \quad \alpha \left(|\Phi^+\rangle^{A'A} + |\Phi^-\rangle^{A'A} \right) |1\rangle^B + \\ & \quad \left. \beta \left(|\Phi^+\rangle^{A'A} - |\Phi^-\rangle^{A'A} \right) |1\rangle^B \right\} \end{aligned}$$

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$$\begin{aligned} &= \frac{1}{2} \left\{ |\Phi^+\rangle^{A'A} (\alpha|0\rangle^B + \beta|1\rangle^B) + \right. \\ &\quad |\Phi^-\rangle^{A'A} (\alpha|0\rangle^B - \beta|1\rangle^B) + \\ &\quad |\Psi^+\rangle^{A'A} (\alpha|1\rangle^B + \beta|0\rangle^B) + \\ &\quad \left. |\Psi^-\rangle^{A'A} (\alpha|1\rangle^B - \beta|0\rangle^B) \right\} \\ &= \frac{1}{2} \left[|\Phi^+\rangle^{A'A} |\psi\rangle^B + |\Phi^-\rangle^{A'A} z |\psi\rangle^B + \right. \\ &\quad \left. |\Psi^+\rangle^{A'A} x |\psi\rangle^B + |\Psi^-\rangle^{A'A} xz |\psi\rangle^B \right] \end{aligned}$$

Now can outline protocol

1) Alice performs Bell measurement on her systems $A'A$. State collapses to

$$\begin{aligned} &|\Phi^+\rangle^{A'A} |\psi\rangle^B, \\ &|\Phi^-\rangle^{A'A} z |\psi\rangle^B, \\ &|\Psi^+\rangle^{A'A} x |\psi\rangle^B, \\ &|\Psi^-\rangle^{A'A} xz |\psi\rangle^B \end{aligned}$$

w/ uniform probability $1/4$

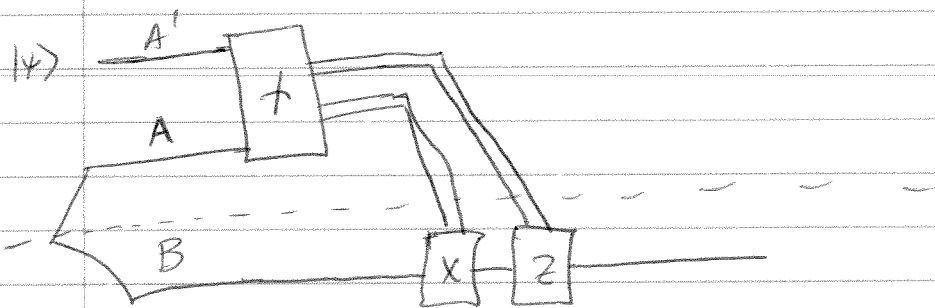
Alice knows which state Bob has.

Bob knows nothing. His state ~~is~~ is π^B .

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- 2) Alice transmits two classical bits to Bob, Bob then knows exactly which state he has
- 3) Bob performs a restoration operation



Note: 1) teleportation process works for any input state (universal)

2) does not violate no-cloning theorem (state is teleported, not copied)

3) not instantaneous

resource inequality is

$$[qq] + 2[c \rightarrow c] \geq [q \rightarrow q]$$

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Optimality of the Three Protocols

$$ED: [q \rightarrow q] \geq [qq]$$

$$SD: [q \rightarrow q] + [qq] \geq 2[c \rightarrow c]$$

$$TP: 2[c \rightarrow c] + [qq] \geq [q \rightarrow q]$$

are these rates optimal?

1st: could there be some protocol where

$$[q \rightarrow q] \geq E [qq]$$

$$\nexists E > 1 ?$$

Suppose it exists.

Assuming free classical comm.,
TP becomes $[qq] \geq [q \rightarrow q]$
can combine w/ TP to get

$$[q \rightarrow q] \geq E [qq] \geq E [q \rightarrow q]$$

can then bootstrap & achieve an unbounded amount of quantum communication, which is impossible (later justify w/ quantum capacity theorem)

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What about optimality of dense coding?

Suppose unlimited amount of entanglement & there is some protocol that gives ~~suppose dense coding generates~~

$$\{q \rightarrow q\} + \infty \{qq\} \geq 2R \{c \rightarrow c\} + \infty \{qq\}$$

where $R > 1$

then can chain this w/ TP:

$$\begin{aligned} \{q \rightarrow q\} + \infty \{qq\} &\geq 2R \{c \rightarrow c\} + \infty \{qq\} \\ &\geq R \{q \rightarrow q\} + \infty \{qq\} \end{aligned}$$

can chain over & over & achieve

$$\{q \rightarrow q\} + \infty \{qq\} \geq R^k \{q \rightarrow q\} + \infty \{qq\}$$

& make arbitrarily ^{high} amount of quantum comm. from just one noiseless qubit & entanglement

violates no-signaling

(can later justify w/ EA capacity theorem)

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Three Noiseless Qudit Protocols

- require this form for use w/
protocols of arbitrary dimension

noiseless qudit channel:

$$|i\rangle^A \rightarrow |i\rangle^B \quad \log_2 d \text{ qudits}$$

noiseless edit channel:

$$\begin{aligned} |i\rangle\langle i|^A &\rightarrow |i\rangle\langle i|^B && \log_2 d \text{ edits} \\ |i\rangle\langle j|^A &\rightarrow 0 && i \neq j \end{aligned}$$

noiseless edit is

$$|\Phi\rangle^{AB} = \frac{1}{\sqrt{d}} \sum_{i=0}^{d-1} |i\rangle^A |i\rangle^B \quad \log_2 d \text{ ebits}$$

Entanglement Distribution

Alice does $CNOT_d \cdot F^A |0\rangle^A |0\rangle^{A'}$

where $F: |e\rangle \rightarrow \frac{1}{\sqrt{d}} \sum_{j=0}^{d-1} \exp\left\{\frac{2\pi i e j}{d}\right\} |j\rangle$

$$CNOT_d = \sum_{j=0}^{d-1} |j\rangle\langle j|^A \otimes X(j)^{A'}$$

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so that

$$F^A |0\rangle^A |0\rangle^{A'} = \frac{1}{\sqrt{d}} \sum_{j=0}^{d-1} |j\rangle^A |0\rangle^{A'}$$

$$\begin{aligned} \xrightarrow{\text{CNOT}} & \sum_{j'=0}^{d-1} (|j'\rangle^{A'} |0\rangle^A \otimes X(j')) \left(\frac{1}{\sqrt{d}} \sum_{j=0}^{d-1} |j\rangle^A |0\rangle^{A'} \right) \\ &= \frac{1}{\sqrt{d}} \sum_{j',j=0}^{d-1} |j'\rangle^{A'} |j\rangle^A \otimes |j\rangle^{A'} \\ &= \frac{1}{\sqrt{d}} \sum_{j=0}^{d-1} |j\rangle^A |j\rangle^{A'} = |\Phi\rangle^{AA'} \end{aligned}$$

noiseless qudit channel

$|\Phi\rangle^{AB}$

resource inequality: $\log d \{q \rightarrow q\} \geq \log d \{qq\}$

Qudit Dense Coding

1) Alice & Bob share

$|\Phi\rangle^{AB}$

Alice applies one of $\left\{ X(x) Z(z) \right\}_{x,z=0}^{d-1}$

where $X(x) |j\rangle = |j \oplus x\rangle$

$Z(z) |j\rangle = \exp\left\{ \frac{i2\pi zj}{d} \right\} |j\rangle$

2) Alice sends ^{her half of} a state $|\Phi_{x,z}\rangle^{AB}$ over qudit channel.

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3) Bob measures in basis

$$\{|\Phi_{x,z}\rangle^{AB}\} \quad (\text{basis is orthonormal})$$

resource inequality is

$$\log d \{q \rightarrow q\} + \log d \{q q\} \geq 2 \log d \{q \rightarrow q\}$$

Quantum Teleportation

global state is

$$|\psi\rangle^{A'} |\Phi\rangle^{AB}$$

$$\text{where } |\psi\rangle^{A'} = \sum_{i=0}^{d-1} \alpha_i |i\rangle^{A'}$$

$$\& \quad |\Phi\rangle^{AB} = \frac{1}{\sqrt{d}} \sum_{j=0}^{d-1} |j\rangle^A |j\rangle^B$$

1) Alice measures $A'A$ in Bell basis

$$\{|\Phi_{i,j}\rangle^{A'A}\}$$

2) transmits $i+j$ to Bob

3) Bob does $Z^B(-j) X^B(-i)$

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Analyze post-measurement state (unnormalized)

$$|\Phi_{ij}\rangle \langle \Phi_{ij}|^{A'A} \quad |\psi\rangle^{RA'} \quad |\Phi\rangle^{AB}$$

projector for $|\Phi_{ij}\rangle$
more general w/ reference system

$$|\Phi_{ij}\rangle^{A'A} \langle \Phi|^{A'A} (U_{ij}^{A'})^\dagger |\psi\rangle^{RA'} |\Phi\rangle^{AB}$$

↓ transpose trick

$$|\Phi_{ij}\rangle^{A'A} \langle \Phi|^{A'A} (U_{ij}^{T A'})^\dagger |\psi\rangle^{RA'} |\Phi\rangle^{AB}$$

↓ commute ↓

$$|\Phi_{ij}\rangle^{A'A} \langle \Phi|^{A'A} |\psi\rangle^{RA'} U_{ij}^{*A} |\Phi\rangle^{AB}$$

↓ transpose trick

$$|\Phi_{ij}\rangle^{A'A} \langle \Phi|^{A'A} |\psi\rangle^{RA'} (U_{ij}^\dagger)^B |\Phi\rangle^{AB}$$

commute ↓

(*)

$$(U_{ij}^\dagger)^B \quad |\Phi_{ij}\rangle^{A'A} \quad \langle \Phi|^{A'A} |\psi\rangle^{RA'} |\Phi\rangle^{AB}$$

evaluate

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$$\langle \Phi | A'A | \Psi \rangle_{RA'} | \Phi \rangle_{AB}$$

$$= \langle \Phi | A'A | \Phi \rangle_{AB} | \Psi \rangle_{RA'}$$

$$= \left(\frac{1}{\sqrt{d}} \sum_{i=0}^{d-1} \langle i | A' \langle i | A \right) \left(\frac{1}{\sqrt{d}} \sum_{j=0}^{d-1} | j \rangle^A | j \rangle^B \right) | \Psi \rangle_{RA'}$$

$$= \left(\frac{1}{d} \sum_{i,j=0}^{d-1} \langle i | A' \langle i | j \rangle^A | j \rangle^B \right) | \Psi \rangle_{RA'}$$

$$= \left[\frac{1}{d} \sum_{i=0}^{d-1} \langle i | A' | i \rangle^B \right] | \Psi \rangle_{RA'}$$

$$= \frac{1}{d} \left[\sum_{i=0}^{d-1} | i \rangle^B \langle i | A' \right] | \Psi \rangle_{RA'}$$

noiseless qudit channel from
 $A' \rightarrow B$

$$= \frac{1}{d} I^{A' \rightarrow B} | \Psi \rangle_{RA'}$$

$$= \frac{1}{d} | \Psi \rangle_{RB}$$

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Back to (*)

$$\frac{1}{d} (U_{ij}^{\dagger})^B |\Phi_{ij}\rangle^{A'A} |\psi\rangle^{RB}$$
$$= \frac{1}{d} |\Phi_{ij}\rangle^{A'A} (U_{ij}^{\dagger})^B |\psi\rangle^{RB} \langle \phi | \phi \rangle^{A'ARB}$$

$$\text{Tr} \left\{ |\phi\rangle\langle\phi|^{A'ARB} \right\} = \frac{1}{d^2}$$

↑
probability uniform &
independent
of state
 $|\psi\rangle^{RB}$

Bob's state w/o classical data is

$$\frac{1}{d^2} \sum_{ij} U_{ij}^B |\psi\rangle\langle\psi|^{RB} (U_{ij}^{\dagger})^B =$$
$$\psi^R \otimes \pi^B$$

w/ classical data, it is

$(U_{ij}^{\dagger})^B |\psi\rangle^{RB}$ ↓ he applies

U_{ij} to bring it to $|\psi\rangle^{RB}$
teleportation!