

Lecture 6

12 FEB 2014

Time to develop the quantum circuit model of computation

Important single-qubit operations

X ≡ [0 1; 1 0], Y ≡ [0 -i; i 0], Z ≡ [1 0; 0 -1]

H ≡ 1/√2 [1 1; 1 -1], S ≡ [1 0; 0 i], T ≡ [1 0; 0 exp(iπ/4)]

↑ phase gate

↑ T gate

can show that H ≡ (X+Z)/√2, S = T^2

useful to define rotation operators

Rx(θ) = e^{-iθX/2} Ry(θ) + Rz(θ)

defined similarly

↑ represents a rotation about x axis in Bloch sphere

will take basic gate set to be {CNOT, H, T}

(2)

First important theorem:

Any single qubit unitary operator can be decomposed as

$$e^{i\delta} \begin{bmatrix} e^{i\alpha/2} & 0 \\ 0 & e^{-i\alpha/2} \end{bmatrix} \begin{bmatrix} \cos \theta/2 & \sin \theta/2 \\ -\sin \theta/2 & \cos \theta/2 \end{bmatrix} \begin{bmatrix} e^{-i\beta/2} & 0 \\ 0 & e^{i\beta/2} \end{bmatrix}$$

plausible: any 2×2 ~~unitary~~ matrix has 8 parameters, but unitarity introduces 4 constraints, leaving 4 parameters

can express any 2×2 unitary as

$$\begin{bmatrix} e^{i(\delta + \alpha/2 + \beta/2)} \cos \theta/2 & e^{i(\delta + \alpha/2 - \beta/2)} \sin \theta/2 \\ -e^{i(\delta - \alpha/2 + \beta/2)} \sin \theta/2 & e^{i(\delta - \alpha/2 - \beta/2)} \cos \theta/2 \end{bmatrix}$$

we then get the factorization above

can rewrite the factorization above as

$$e^{i\delta} R_z(\alpha) R_y(\theta) R_z(\beta)$$

③

Useful corollary:

Let U be a single-qubit unitary gate.

Then $\exists A, B, C$ (all unitary) such

$$\text{that } ABC = I \quad \dagger$$

$$U = e^{i\phi} A X B X C$$

$$\text{Take } A = R_z(\alpha) R_y(\theta/2)$$

$$B = R_y(-\frac{\theta}{2}) R_z(-(\frac{\beta+\alpha}{2}))$$

$$C = R_z(\frac{\beta-\alpha}{2})$$

$$\text{So } ABC = I \quad (\text{by inspection})$$

$$\underline{A X B X C}$$

$$X B X = X R_y(-\frac{\theta}{2}) X X R_z(-(\frac{\beta+\alpha}{2})) X$$

$$= R_y(\frac{\theta}{2}) R_z(\frac{\beta+\alpha}{2})$$

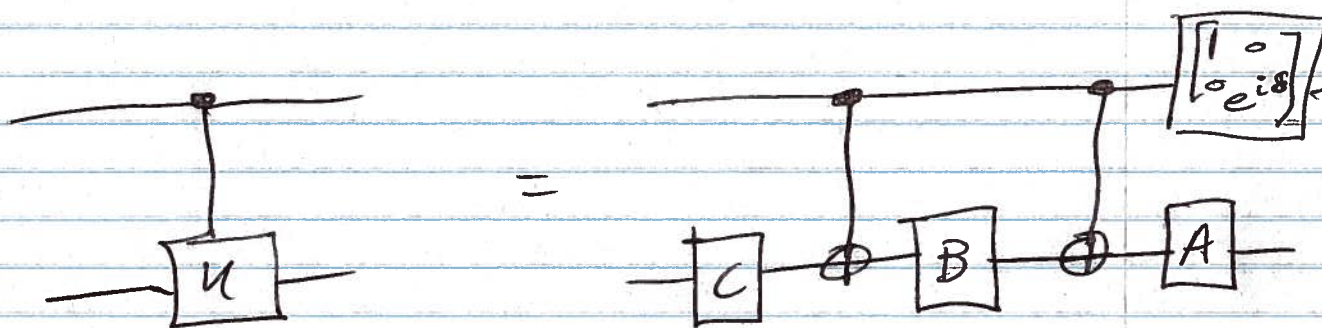
$$\begin{aligned} \text{then } & \underbrace{R_z(\alpha) R_y(\frac{\theta}{2}) R_y(\frac{\theta}{2}) R_z(\frac{\beta+\alpha}{2})}_{A} \underbrace{R_z(\frac{\beta-\alpha}{2})}_C \\ & = R_z(\alpha) R_y(\theta) R_z(\beta) \end{aligned}$$

(4)

utility of this decomposition is in promoting a single qubit unitary to a controlled one:

$$|0\rangle\langle 0| \otimes I + |1\rangle\langle 1| \otimes U$$

$$|ij\rangle \rightarrow (I \otimes U^i) |ij\rangle$$



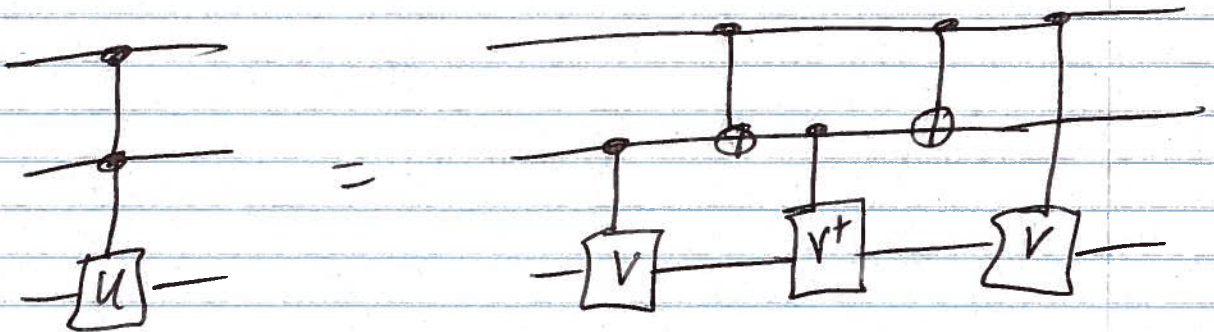
so if we have CNOTs + arbitrary single-qubit unitaries, then we can implement an arbitrary ^{2-qubit} controlled-unitary

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controlling on multiple qubits

$$C^n(U) |x_1 \dots x_n\rangle |\psi\rangle = |x_1 \dots x_n\rangle U^{x_1 \dots x_n} |\psi\rangle$$

example:



where $V^2 = U$

need Toffoli gate as well

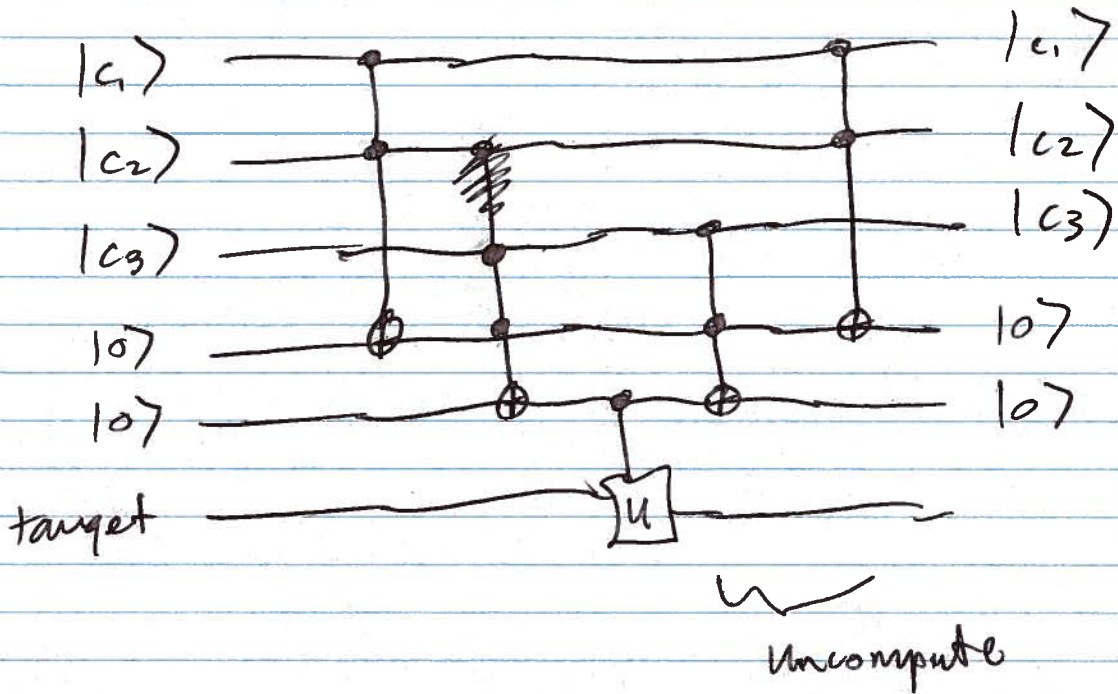


controlled-controlled-NOT

rather complicated implementation

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To implement controlled- U , we do



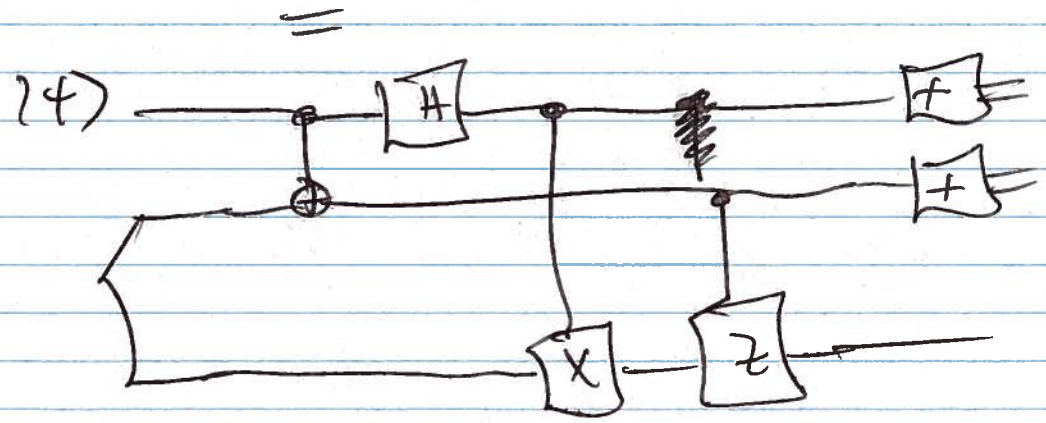
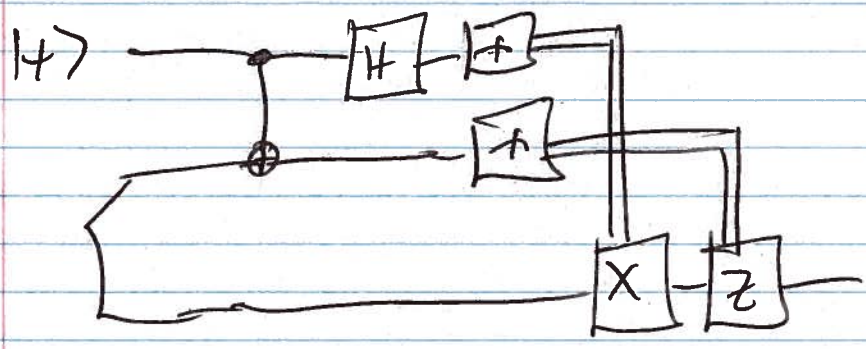
only linear overhead

regarding measurements, we will only mention the "principle of deferred measurement"

can always move measurements to the end of the computation w/o changing the operation of circuit. replace conditional operations w/ controlled operations

\Rightarrow adaptive strategies do not increase computational power

Example: teleportation



universality of quantum gates:

would like to prove that a discrete gate set can simulate any n-qubit unitary to any desired accuracy:

For $\forall U \in \mathcal{U} \ \forall \epsilon > 0 \ \exists V_1, \dots, V_M$ such that

$$\max_{|\psi\rangle} \| U|\psi\rangle - V_M \dots V_1 |\psi\rangle \|_2 \leq \epsilon$$

Then the ~~circuit~~ unitary of circuit will be indistinguishable up to an ϵ -error

Important question: What is the overhead in the simulation?

Begin by showing that any unitary on n qubits can be implemented by CNOTs & single-qubit unitaries.

1st understand how to decompose unitaries using two-level unitaries

$$U = \begin{bmatrix} a & d & g \\ b & e & h \\ c & f & j \end{bmatrix}$$

want to find 2-level unitaries such that

$$U_3 U_2 U_1 U = I$$

if $b=0$ then $U_1 = I$

$$\text{if } b \neq 0 \text{ then } U_1 = \frac{1}{\|(a,b)\|_2} \begin{bmatrix} a^* & b^* & 0 \\ b & -a & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$\text{So } U_1 U = \begin{bmatrix} a' & d' & g' \\ 0 & e' & h' \\ c' & f' & j' \end{bmatrix}$$

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Now ~~set~~ if $c' = 0$ set

$$U_2 = \begin{bmatrix} a'^* & & \\ & 1 & \\ & & 1 \end{bmatrix}$$

if $c' \neq 0$ set

$$U_2 = \frac{1}{\|(a', c')\|_2} \begin{bmatrix} a'^* & & c'^* \\ & 1 & \\ c' & & -a' \end{bmatrix}$$

then $U_2 U_1 U =$

$$\begin{bmatrix} 1 & d'' & g'' \\ 0 & e'' & h'' \\ 0 & f'' & j'' \end{bmatrix}$$

From unitarity, $d'', g'' = 0$

So then set $U_3 = \begin{bmatrix} 1 & & \\ & e'' & h'' \\ & f'' & j'' \end{bmatrix}$

$e''^* f''^*$
 $h''^* j''^*$

can do a similar kind of thing for larger dimensional unitaries

Now show that we can use CNOTs ~~as well~~ & single-qubit unitaries

Before we showed that we can decompose U in terms of two-level matrices that act nontrivially on a 2D subspace & trivially on ~~the~~ complementary subspace.

Suppose the basis for the subspace is $\{|s\rangle, |t\rangle\}$ where $|s\rangle = |s_1 \dots s_n\rangle$
 $|t\rangle = |t_1 \dots t_n\rangle$

We use the classical idea of Gray codes to effect each two-level transformation.

Example: Suppose 2-level unitary is

$$\begin{bmatrix} a & & & & & c \\ & 1 & & & & \\ & & 1 & & & \\ & & & 1 & & \\ & & & & 1 & \\ b & & & & & d \end{bmatrix}$$

The unitary acts nontrivially on the space ~~spanned by~~ $\text{span}\{|000\rangle, |111\rangle\}$

So we find a Gray code connecting these states

would like a way to

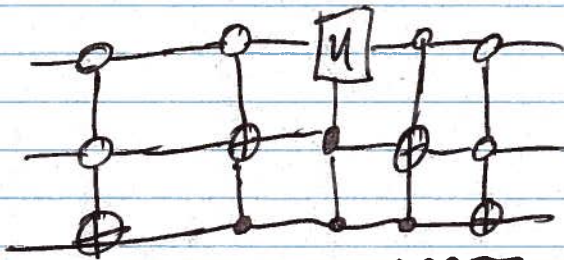
~~change~~ permute this basis to $\{|011\rangle, |111\rangle\}$ then act w/

000
001
011
111

at most 1 bit changes in each transition

to effect transformation

so we perform



controlled on 00, flip the third takes

000 → 001 → 011

since the Toffoli can be realized w/ single-qubit unitaries + CNOTs, ^(and same for C-U) we have achieved the goal.