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Strong converse exponents for
a quantum channel discrimination
problem

(joint w/ Cooney &
Mason) in
arXiv:1408.3373)

Umegaki ('62)

Relative entropy is a fundamental
quantity in QIT, defined as

$$D(\rho \parallel \sigma) = \text{Tr} \left\{ \rho [\log \rho - \log \sigma] \right\}$$

Interesting that there are other
generalizations of classical
relative entropy such as Belavukha-
Staszewski

$$D_{BS}(\rho \parallel \sigma) = \text{Tr} \left\{ \rho \log(\rho^{1/2} \sigma^{-1} \rho^{1/2}) \right\}$$

and in fact many other ways
to generalize...

What sets the Umegaki one
apart is that it has an appealing
operational interpretation in
terms of the quantum Stein's lemma
tho. those don't...

Setting of quantum Stein's lemma is asymmetric quantum hypothesis testing.

What is hypothesis testing?

Someone prepares a quantum system in the state ρ or σ , but does not tell you which.

You then get to make a measurement $\{Q, I-Q\}$ on system.

\uparrow identity \uparrow identity σ
 ρ

There are two kinds of errors you can make:

(Type I) system prepared as ρ
but measured as σ

$$\alpha(Q) \equiv \text{Tr} \{ (I-Q) \rho \}$$

(Type II) system prepared as σ
but measured as ρ

$$\beta(Q) \equiv \text{Tr} \{ Q \sigma \}$$

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There is a non-trivial tradeoff
between the two kinds of errors

$$\alpha(Q) + \beta(Q) > 0 \quad \text{unless } \rho \perp \sigma$$

asymmetric setting explores the
trade-off by imposing a constraint
on Type I error & minimizing
Type II error subject to the constraint

Define

$$P_{\epsilon}(p \parallel \sigma) = \min \left\{ \beta(Q) : 0 \leq Q \leq I \wedge \alpha(Q) \leq \epsilon \right\}$$

can be useful to define the

"hypothesis testing relative entropy" as

$$D_{\#}^{\epsilon}(p \parallel \sigma) = -\log P_{\epsilon}(p \parallel \sigma)$$

obeys monotonicity under quantum operations

$$D_{\#}^{\epsilon}(p \parallel \sigma) \geq D_{\#}^{\epsilon}(N(p) \parallel N(\sigma))$$

b/c performing the optimal test can
always do better than doing a channel +
then a test

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quantum Stein's lemma applies
to the case of distinguishing
 $\rho^{\otimes n}$ from $\sigma^{\otimes n}$ using
a collective measurement $\{Q_n, I^{\otimes n} - Q_n\}$

Direct part of theorem: (Hiai-Petz '91)

\exists a sequence of measurements $\{Q_n^*, \dots\}$
such that for all $\epsilon > 0$ + suff. large n
$$\alpha_n(Q_n) \leq \epsilon$$

while
$$\beta_n(Q_n) \leq e^{-n[D(\rho||\sigma) - \delta]}$$

where $\delta > 0$

This gives the inequality

$$\lim_{n \rightarrow \infty} \frac{1}{n} D_H^\epsilon(\rho^{\otimes n} || \sigma^{\otimes n}) \geq D(\rho || \sigma)$$

Strong converse part of theorem:

For any sequence of measurements such that $\{Q_n\}$

$$P_n(Q_n) \leq e^{-n[D(\rho||\sigma) + \delta]}$$

w/ $\delta > 0$

then $\alpha_n(Q_n) \rightarrow 1$

so that there is ~~no~~ trade-off

equivalent statement is of strong converse

$$\lim_{n \rightarrow \infty} \frac{1}{n} D_H^\epsilon(\rho^{\otimes n} || \sigma^{\otimes n}) \leq D(\rho || \sigma)$$

can prove it somewhat easily via the Renyi relative entropy & the following lemma

$$D_\alpha(\rho || \sigma) = \frac{1}{\alpha - 1} \log \text{Tr} \{ \rho^\alpha \sigma^{1-\alpha} \}$$

$$D_H^\epsilon(\rho || \sigma) \leq D_\alpha(\rho || \sigma) + \frac{\alpha}{\alpha - 1} \log \left(\frac{1}{1 - \epsilon} \right)$$

for $\alpha > 1$ & $\epsilon \in (0, 1)$

How to prove: choose $\{Q, I-Q\}$ to be optimum for D_H^ϵ

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Apply to tensor-power states & regularize:

$$\begin{aligned} \frac{1}{n} D_H^\epsilon(\rho^{\otimes n} \| \sigma^{\otimes n}) &\leq \frac{1}{n} D_\alpha(\rho^{\otimes n} \| \sigma^{\otimes n}) \\ &\quad + \frac{1}{n} \left(\frac{\alpha}{\alpha-1} \right) \log \left(\frac{1}{1-\epsilon} \right) \\ &= D_\alpha(\rho \| \sigma) + \frac{1}{n} (\dots) \end{aligned}$$

Take limit as $n \rightarrow \infty$

to get

$$\lim_{n \rightarrow \infty} \frac{1}{n} D_H^\epsilon(\rho^{\otimes n} \| \sigma^{\otimes n}) \leq D_\alpha(\rho \| \sigma)$$

but holds for every α , so

take limit as ~~α~~ $\alpha \rightarrow 1$ to

get result since

$$\lim_{\alpha \rightarrow 1} D_\alpha(\rho \| \sigma) = D(\rho \| \sigma)$$

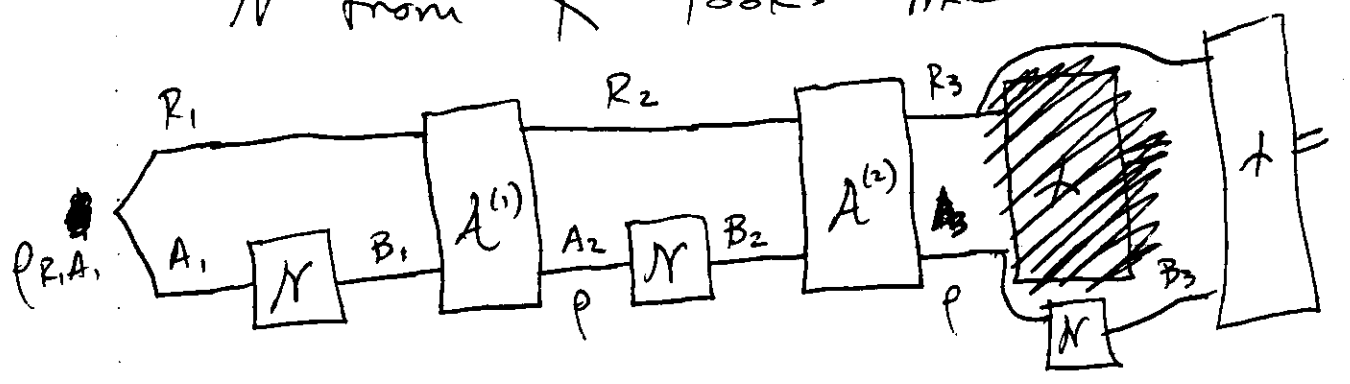
New

Setting we are interested in is

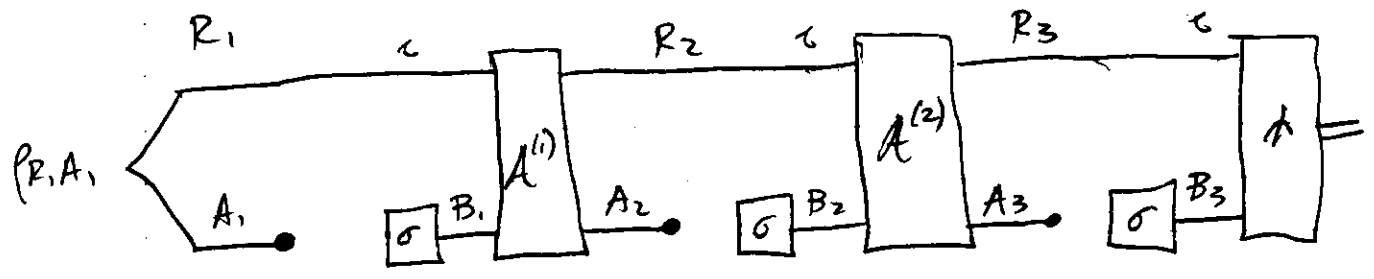
adaptive channel discrimination
between N + receiver

$$\text{channel } \mathcal{P}(x) \equiv \text{Tr}\{X\} \cdot \sigma$$

General strategy for discriminating
 N from \mathcal{R} looks like



vs.



We would like to establish a Stein's lemma in this setting

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Direct part is easy:

Just use the old Stein's lemma

w/ $\phi_{RA}^{\otimes n}$ & no adaptation

if channel is \mathcal{N} , then output
is $[\mathcal{N}_{A \rightarrow B}(\phi_{RA})]^{\otimes n}$

if channel is \mathcal{R} , then output
is $[\phi_R \otimes \sigma_B]^{\otimes n}$

& optimal decay exponent is

$$\max_{\phi_{RA}} D(\mathcal{N}_{A \rightarrow B}(\phi_{RA}) \parallel \phi_R \otimes \sigma_B)$$

↑ can optimize...

Strong converse part:

use "sandwiched" Rényi relative entropy

from arXiv: 1306.3142 &

arXiv: 1306.1586

$$\tilde{D}_\alpha(\rho \parallel \sigma) \equiv \frac{1}{\alpha-1} \log \text{Tr} \left\{ \left(\sigma^{\frac{1-\alpha}{2\alpha}} \rho \sigma^{\frac{1-\alpha}{2\alpha}} \right)^\alpha \right\}$$

$$= \frac{\alpha}{\alpha-1} \log \left\| \sigma^{\frac{1-\alpha}{2\alpha}} \rho \sigma^{\frac{1-\alpha}{2\alpha}} \right\|_\alpha$$

w/ $\|A\|_\alpha \equiv \left(\text{Tr} \{ |A|^\alpha \} \right)^{1/\alpha}$

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use lemma

$$D_H^\epsilon(\rho \parallel \sigma) \leq \tilde{D}_\alpha(\rho \parallel \sigma) + \left(\frac{\alpha}{\alpha-1}\right) \log\left(\frac{1}{1-\epsilon}\right)$$

Apply this lemma to the last step of picture

$$D_H^\epsilon(N_{A_3 \rightarrow B_3}(\rho_{R_3 A_3}) \parallel \tau_{R_3} \otimes \sigma_{B_3}) \leq \tilde{D}_\alpha(N_{A_3 \rightarrow B_3}(\rho_{R_3 A_3}) \parallel \tau_{R_3} \otimes \sigma_{B_3}) + \left(\frac{\alpha}{\alpha-1}\right) \log\left(\frac{1}{1-\epsilon}\right) \quad (\Psi)$$

focus on \downarrow rewrite as

$$\begin{aligned} & \frac{\alpha}{\alpha-1} \log \left\| (\tau_{R_3} \otimes \sigma_{B_3})^{\frac{1-\alpha}{2\alpha}} N_{A_3 \rightarrow B_3}(\rho_{R_3 A_3}) (\tau_{R_3} \otimes \sigma_{B_3})^{\frac{1-\alpha}{2\alpha}} \right\|_\alpha \\ &= \frac{\alpha}{\alpha-1} \log \left\| \sigma_{B_3}^{\frac{1-\alpha}{2\alpha}} N_{A_3 \rightarrow B_3} \left(\tau_{R_3}^{\frac{1-\alpha}{2\alpha}} \rho_{R_3 A_3} \tau_{R_3}^{\frac{1-\alpha}{2\alpha}} \right) \sigma_{B_3}^{\frac{1-\alpha}{2\alpha}} \right\|_\alpha \\ &= \frac{\alpha}{\alpha-1} \log \left\| \left(\oplus_{\sigma} N_{A_3 \rightarrow B_3} \right) \left(\tau_{R_3}^{\frac{1-\alpha}{2\alpha}} \rho_{R_3 A_3} \tau_{R_3}^{\frac{1-\alpha}{2\alpha}} \right) \right\|_\alpha \end{aligned}$$

where $\oplus_{\sigma}(X) \equiv \sigma_{B_3}^{\frac{1-\alpha}{2\alpha}}(X) \sigma_{B_3}^{\frac{1-\alpha}{2\alpha}}$

Focus on expression in log & rewrite as

$$\begin{aligned}
 & \frac{\| (\oplus_{\sigma} \circ X_{A_3 \rightarrow B_3}) (\tau_{R_3}^{\frac{1-\alpha}{2\alpha}} \rho_{R_3 A_3} \tau_{R_3}^{\frac{1-\alpha}{2\alpha}}) \|_{\alpha}}{\| \tau_{R_3}^{\frac{1-\alpha}{2\alpha}} \rho_{R_3} \tau_{R_3}^{\frac{1-\alpha}{2\alpha}} \|_{\alpha}} \cdot \| \tau_{R_3}^{\frac{1-\alpha}{2\alpha}} \rho_{R_3} \tau_{R_3}^{\frac{1-\alpha}{2\alpha}} \|_{\alpha} \\
 \leq & \left[\sup_{X_{R_3 A_3} > 0} \frac{\| (\oplus_{\sigma} \circ X_{A_3 \rightarrow B_3}) (X_{R_3 A_3}) \|_{\alpha}}{\| X_{R_3} \|_{\alpha}} \right] \cdot \| \tau_{R_3}^{\frac{1-\alpha}{2\alpha}} \rho_{R_3} \tau_{R_3}^{\frac{1-\alpha}{2\alpha}} \|_{\alpha} \\
 = & \| \oplus_{\sigma} \circ X \|_{\alpha, 1 \rightarrow \alpha} \cdot \| \tau_{R_3}^{\frac{1-\alpha}{2\alpha}} \rho_{R_3} \tau_{R_3}^{\frac{1-\alpha}{2\alpha}} \|_{\alpha}
 \end{aligned}$$

plug back into log to find that (*)

$$\begin{aligned}
 & \leq \frac{\alpha}{\alpha-1} \log \| \oplus_{\sigma} \circ X \|_{\alpha, 1 \rightarrow \alpha} + \frac{\alpha}{\alpha-1} \log \| \tau_{R_3}^{\frac{1-\alpha}{2\alpha}} \rho_{R_3} \tau_{R_3}^{\frac{1-\alpha}{2\alpha}} \|_{\alpha} \\
 & = \quad \quad \quad + \tilde{D}_{\alpha}(\rho_{R_3} \| \tau_{R_3}) \\
 & \leq \quad \quad \quad + \tilde{D}_{\alpha}(X_{A_2 \rightarrow B_2}(\rho_{R_2 A_2}) \| \tau_{R_2} \circ \sigma_{B_2}) \\
 & \uparrow \\
 & \text{monotonicity of } \tilde{D}_{\alpha} \text{ under } \text{Tr}_{A_3} \circ A^{(\alpha)} \quad \forall \alpha > 1
 \end{aligned}$$

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CB, $1 \rightarrow \alpha$ norm defined as
of M

$$\|M\|_{CB, 1 \rightarrow \alpha} = \max_{P_R} \left\| \left(P_R^{1/\alpha} \otimes I_B \right) \Gamma_{RB}^M \left(P_R \otimes P_B \right)^{1/\alpha} \right\|_{\alpha}$$

where $\Gamma_{RB}^M \equiv M_{A \rightarrow B}(\Gamma_{RA})$

$$\Gamma_{RA} \equiv \sum_i |i\rangle_R |i\rangle_A$$

(1)

Iterate this "n times" to get

$$\begin{aligned} D_H^\epsilon & \left(N_{A_n \rightarrow B_n}(\rho_{R_n A_n}) \parallel \tau_{R_n} \otimes \sigma_{B_n} \right) \\ & \leq n \cdot \frac{2}{2-1} \log \left\| \oplus_{\sigma} N \right\|_{CB, 1 \rightarrow 2} + \left(\frac{2}{2-1} \right) \log \left(\frac{1}{1-\epsilon} \right) \\ & = n \cdot \sup_{\phi_{RA}} \tilde{D}_2 \left(N_{A \rightarrow B}(\phi_{RA}) \parallel \phi_R \otimes \sigma_B \right) \end{aligned}$$

~~Divide~~ Divide by n & take limit
as $n \rightarrow \infty$ & then as $2 \rightarrow 1$
to get upper bound of

$$\sup_{\phi_{RA}} D \left(N_{A \rightarrow B}(\phi_{RA}) \parallel \phi_R \otimes \sigma_B \right)$$

applications are in quantum illumination
to demonstrate that tensor-power
strategies are optimal

can also show strong converse
for quantum feedback assisted communication